



The key to FPGA-ASIC design

Creators of the  
Clash compiler



The key to  
FPGA design

[qbaylogic.com](http://qbaylogic.com)

# ghc-typelits-proof-assist

## A Haskell Plugin for Type-Nat Proofs

Felix Klein

HIW – June 6th, 2025



The key to FPGA-ASIC design

# Problem

```
import GHC.TypeNats

constrainedOp :: n ≤ m ⇒ SNat n → SNat m → ()
constrainedOp _ _ = ()

cOpApp :: n + 1 ≤ m ⇒ SNat n → SNat m → ()
cOpApp = constrainedOp
```

```
<interactive>:13:11: error: [GHC-64725]
```

- Cannot satisfy:  $n \leq m$
- In the expression: `constrainedOp`
- In an equation for ‘`cOpApp`’: `cOpApp = constrainedOp`

## Existing (partial) Solutions

- `ghc-typelits-natnormalise`
- `ghc-typelits-extra`
- `ghc-typelits-presburger`
- `ghc-typelits-sopsat`

They work well for “*simple*” expressions, but reach their limits as soon as the required deductions get too complex.

**Any alternative solutions?** 🐈 “unsafeCoerce” 🐈

**Idea of this plugin:** Use an external proof assistant to create the required evidence.

## How does it work

```
{-# LANGUAGE UndecidableInstances #-}  
{-# LANGUAGE UndecidableSuperClasses #-}  
  
...  
  
instance  
  ( n + 1 ≤ m  
  ) ⇒ Lemma n m  
class -- =>  
  ( n ≤ m  
  ) ⇒ Lemma n m  
{-/ Proof (Coq): Lemma /-}
```

1. Creates “Lemma.v”, a Coq proof file with the signature of the proof.

## How does it work

```
Require Import Coq.Init.Nat.  
Require Import Coq.Init.Peano.  
  
Require Import Coq.Arith.PeanoNat.  
  
Lemma hsLemma: forall n m, n + 1 <= m -> n <= m.
```

2. Develop some proof in your favorite Coq IDE.

## How does it work

```
...  
  
Lemma hsLemma: forall n m, n + 1 <= m -> n <= m.  
  intros.  
  apply le_S in H.  
  apply le_S_n.  
  rewrite <- PeanoNat.Nat.add_0_1 at 1.  
  rewrite PeanoNat.Nat.add_comm.  
  rewrite PeanoNat.Nat.add_succ_1.  
  rewrite PeanoNat.Nat.add_comm.  
  rewrite <- PeanoNat.Nat.add_succ_1.  
  rewrite PeanoNat.Nat.add_comm.  
  apply H.  
Qed.
```

3. Copy back the proof into the Haskell “*proof comment*”.

## How does it work

```
...  
  
instance  
  ( n + 1 ≤ m  
  ) ⇒ Lemma n m  
class -- =>  
  ( n ≤ m  
  ) ⇒ Lemma n m  
{-/ Proof (Coq): Lemma  
  intros.  
  apply le_S in H.  
  apply le_S_n.  
  rewrite <- PeanoNat.Nat.add_0_1 at 1.  
  rewrite PeanoNat.Nat.add_comm.  
  rewrite PeanoNat.Nat.add_succ_1.  
  rewrite PeanoNat.Nat.add_comm.  
  rewrite <- PeanoNat.Nat.add_succ_1.  
  rewrite PeanoNat.Nat.add_comm.  
  apply H.  
/-}
```

## The QED Class

```
data Rewrite (c :: Constraint) = c => Rewrite

class c => QED (c :: Constraint) where
  apply :: (c => a) -> a
  apply x = x

using :: Rewrite c
using = Rewrite
```

⇒ Enables the usage and distribution of the created evidence.

# The QED Class

```
...  
  
instance  
  ( n + 1 ≤ m  
  ) ⇒ Lemma n m  
class -- =>  
  ( n ≤ m  
  ) ⇒ Lemma n m  
{-/ Proof (Coq): Lemma  
  intros.  
  apply le_S in H.  
  apply le_S_n.  
  rewrite <- PeanoNat.Nat.add_0_1 at 1.  
  rewrite PeanoNat.Nat.add_comm.  
  rewrite PeanoNat.Nat.add_succ_1.  
  rewrite PeanoNat.Nat.add_comm.  
  rewrite <- PeanoNat.Nat.add_succ_1.  
  rewrite PeanoNat.Nat.add_comm.  
  apply H.  
-/}  
instance Lemma n m ⇒ QED (Lemma n m)
```

## Using the Evidence

```
-- using 'Rewrite'
cOpApp :: forall n m. n + 1 ≤ m ⇒ SNat n → SNat m → ()
cOpApp | Rewrite ← using @(Lemma n m) = constrainedOp

-- using 'apply'
cOpApp :: forall n m. n + 1 ≤ m ⇒ SNat n → SNat m → ()
cOpApp = apply @(Lemma n m) constrainedOp

-- using the AutoProve feature of the plugin (limited)
usingAuto :: n + 1 ≤ m ⇒ SNat n → SNat m → ()
usingAuto = constrainedOp
```

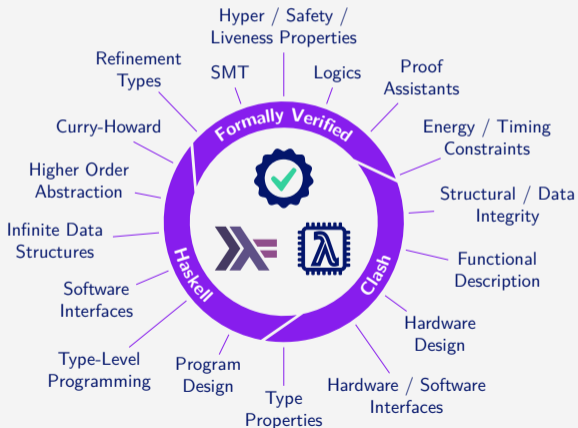
## Next Steps

- Custom type family support
- Support more proof assistants (currently: Agda, Coq)
- GHC 9.12 / 9.14 support (currently: only 9.10)
- Integrate with the other existing plugins



# Clash Formal

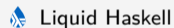
Ecosystem formally verifiable IT



## Proof Assistants



## SMT / Refinement Types



## Specification Frameworks



<https://clash-formal.org>